# ON MINIMIZING QUADRATIC FUNCTIONS WITH SEPARABLE CONVEX CONSTRAINTS 

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#### Abstract

A new active set algorithm for minimizing quadratic functions with separable convex constraints combines the conjugate gradient method with gradient projections. It generalizes recently developed algorithms of quadratic programming constrained by simple bounds. A linear convergence rate in terms of the Hessian spectral condition number is proved. Numerical experiments including frictional 3D contact problems of linear elasticity illustrate the computational performance.


Keywords: Quadratic function, Separable convex constraints, Active set, Conjugate gradient method, Projected gradient, Convergence rate

## 1. Introduction

We shall be concerned with solving

$$
\begin{equation*}
\min _{x \in \Omega} f(x), \tag{1}
\end{equation*}
$$

where $f(x)=\frac{1}{2} x^{\top} A x-x^{\top} b, A \in \mathbb{R}^{n \times n}$ is symmetric positive definite, $b \in \mathbb{R}^{n}, \Omega=\Omega_{1} \times \cdots \times \Omega_{m}$, and $\Omega_{i}=\left\{\boldsymbol{x}_{i} \in \mathbb{R}^{n_{i}}: f_{i}\left(\boldsymbol{x}_{i}\right) \leq 0\right\}$ are defined by continuously differentiable convex functions $f_{i}: \mathbb{R}^{n_{i}} \mapsto \mathbb{R}$ so that $n_{i} \geq 1$, $\sum_{i=1}^{m} n_{i}=n$. Let us note that the feasible set $\Omega$ is separable in the sense that each part $\boldsymbol{x}_{i}$ of $x=\left(\boldsymbol{x}_{1}^{\top}, \ldots, \boldsymbol{x}_{m}^{\top}\right)^{\top}$ is subject to one constraint $\boldsymbol{x}_{i} \in \Omega_{i}$.

This problem includes several independently investigated subproblems originating for instance in duality based methods for the solution of contact problems of linear elasticity:

- If $n_{i}=1$ and $f_{i}\left(x_{i}\right) \equiv l_{i}-x_{i}$ with $l_{i}$ given, we obtain the simple bound

$$
\begin{equation*}
l_{i} \leq x_{i} \tag{2}
\end{equation*}
$$

arising from 2D contact problems [1].

- If $n_{i}=2, \boldsymbol{x}_{i}=\left(x_{2 i-1}, x_{2 i}\right)^{\top}$ and $f_{i}\left(\boldsymbol{x}_{i}\right) \equiv x_{2 i-1}^{2}+x_{2 i}^{2}-r_{i}^{2}$ with $r_{i}$ given, we arrive at

$$
\begin{equation*}
x_{2 i-1}^{2}+x_{2 i}^{2} \leq r_{i}^{2} \tag{3}
\end{equation*}
$$

that can be interpreted as the circular constraint. A source of such constraints is a friction law for 3 D contact problems [3].

## 2. Algorithm

Let $\mathcal{M}=\{1, \ldots, m\}$. The gradient of $f$ at a point $x \in \Omega$ is $g:=g(x)=A x-b$ and the active-set is defined by $\mathcal{A}:=\mathcal{A}(x)=\left\{i \in \mathcal{M} \mid f_{i}\left(\boldsymbol{x}_{i}\right)=0\right\}$. Using the projection $P_{\Omega}: \mathbb{R}^{n} \mapsto \Omega$, we define the projected gradient for a fixed $\widetilde{\alpha} \geq 0$ as

$$
\widetilde{g}:=\widetilde{g}(x)=\frac{1}{\widetilde{\alpha}}(x-P(x-\widetilde{\alpha} g(x)))
$$

The projected gradient characterizes the solution $x^{*}$ of (1) by $\widetilde{g}\left(x^{*}\right)=0$. Our algorithm is based on the fact that the non-zero components of $\widetilde{g}(x)$ at $x \neq x^{*}$ determine the descent directions changing appropriately the activeset. To this end, we introduce components of $\widetilde{\sim}(x)$ called the projected free gradient $\widetilde{\phi}:=\widetilde{\phi}(x)$ and the projected boundary gradient $\widetilde{\beta}:=\widetilde{\beta}(x)$, respectively, as follows:

$$
\begin{array}{ll}
\widetilde{\phi}_{\mathcal{A}}=0, & \widetilde{\phi}_{\mathcal{M} \backslash \mathcal{A}}=\widetilde{g}_{\mathcal{M} \backslash \mathcal{A}} \\
\widetilde{\beta}_{\mathcal{A}}=\widetilde{g}_{\mathcal{A}}, & \widetilde{\beta}_{\mathcal{M} \backslash \mathcal{A}}=0
\end{array}
$$

We combine three steps to generate a sequence $\left\{x^{(l)}\right\}$ that approximates the solution to (1):

- the expansion step: $x^{(l+1)}=x^{(l)}-\widetilde{\alpha} \widetilde{\phi}\left(x^{(l)}\right)$,
- the proportioning step: $x^{(l+1)}=x^{(l)}-\widetilde{\alpha} \widetilde{\beta}\left(x^{(l)}\right)$,
- the conjugate gradient step: $x^{(l+1)}=x^{(l)}-\alpha_{c g}^{(l)} p^{(l)}$, where $\alpha_{c g}^{(l)}$ and the conjugate gradient directions $p^{(l)}$ are computed recurrently; the recurrence starts from $x^{(s)}$ generated by the last expansion or the proportioning step and satisfies $\mathcal{A}\left(x^{(l+1)}\right)=\mathcal{A}\left(x^{(s)}\right)$.

The expansion step may add indices while the proportioning step may release indices to/from the current active-set. The conjugate gradient steps are used to carry out efficiently the minimization of the objective $f$ in the interior of the face $W\left(x^{(s)}\right)=\left\{x \in \Omega \mid x_{\mathcal{A}}:=x_{\mathcal{A}}^{(s)}, \mathcal{A}=\mathcal{A}\left(x^{(s)}\right)\right\}$. Moreover, the algorithm exploits a given constant $\gamma>0$ and the releasing criterion

$$
\begin{equation*}
\widetilde{\beta}\left(x^{(l)}\right)^{\top} g\left(x^{(l)}\right) \leq \gamma \widetilde{\phi}\left(x^{(l)}\right)^{\top} g\left(x^{(l)}\right) \tag{4}
\end{equation*}
$$

to decide which of the steps will be performed.
Algorithm QPC [4] Let $x^{(0)} \in \Omega, \gamma>0, \widetilde{\alpha} \in\left(0,\|A\|^{-1}\right.$ ] and $\epsilon>0$ be given. For $x^{(l)}, x^{(s)}$ known, $0 \leq s \leq l$, where $x^{(s)}$ is computed by the last step expansion or proportioning, choose $x^{(l+1)}$ by the following rules:
(i) If $\left\|\widetilde{g}\left(x^{(l)}\right)\right\| \leq \epsilon$, return $x:=x^{(l)}$.
(ii) If $x^{(l)}$ fulfils (4), try to generate $x^{(l+1)}$ by the conjugate gradient step. If $x^{(l+1)} \in \operatorname{Int} W\left(x^{(s)}\right)$, accept it, otherwise generate $x^{(l+1)}$ by the expansion step.
(iii) If $x^{(l)}$ does not fulfil (4), generate $x^{(l+1)}$ by the proportioning step.

Contrary to simple bound problems analyzed in [2], the algorithm does not exhibit the finite terminating property while the same convergence rate is achieved.

Theorem 1 [5] Let $x^{*} \in \Omega$ be the solution to (1), $\alpha_{\min }$ denote the smallest eigenvalue of $A$ and $\widehat{\gamma}=\max \left\{\gamma, \gamma^{-1}\right\}$. Let $\left\{x^{(l)}\right\}$ be the sequence generated by Algorithm 2.1 with $\epsilon=0$. Then

$$
f\left(x^{(l+1)}\right)-f\left(x^{*}\right) \leq \eta\left(f\left(x^{(l)}\right)-f\left(x^{*}\right)\right)
$$

where

$$
\eta=1-\frac{\widetilde{\alpha} \alpha_{\min }}{2+2 \widehat{\gamma}}<1
$$

The error in the $A$-energy norm is bounded by

$$
\left\|x^{(l)}-x^{*}\right\|_{A}^{2} \leq 2 \eta^{k}\left(f\left(x^{(0)}\right)-f\left(x^{*}\right)\right)
$$

Theorem 1 yields the optimal value of $\eta$ for $\gamma=\hat{\gamma}=1$ and $\widetilde{\alpha}=\|A\|^{-1}$, when

$$
\eta=1-\frac{1}{4} \kappa(A)^{-1}
$$

where $\kappa(A)$ is the spectral condition number of $A$.

## 3. Numerical experiments

We consider a steel brick supported by a rigid foundation that occupies the bounded domain. The boundary is split into three nonempty disjoint parts with different boundary conditions: zero displacements, surface tractions and contact conditions. The contact conditions are represented by the nonpenetration, an effect of friction and the transmission of contact stresses; see [3].

The discrete contact problem with Tresca friction reduces to (1). As the more realistic Coulomb friction law leads to a sequence of Tresca friction problems we can repeatedly apply Algorithm QPC to solve it.

The table summarizes CPU time (in seconds), the number of successive approximations (iter) and the total complexity in terms of matrix-vector
multiplications $\left(n_{A}\right)$ for various primal and dual dofs $3 n_{c}$ and $n=3 m_{c}$, respectively, and for two coefficients of friction $F$. The obtained results are promising, especially, $n_{A}$ is only mildly dependent on the finite element discretization so that the relative complexity $n_{A} / n$ considerably decreases for finer grids.

| dofs |  | $F=0.3$ |  |  |  |  | $F=0.6$ |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
| $3 n_{c}$ | $3 m_{c}$ | time | iter | $n_{A}$ | $n_{A} / n$ | time | iter | $n_{A}$ | $n_{A} / n$ |  |
| 900 | 180 | 4 | 5 | $\mathbf{5 3 5}$ | 2.97 | 6 | 7 | $\mathbf{8 0 1}$ | 4.45 |  |
| 2646 | 378 | 24 | 5 | $\mathbf{6 3 8}$ | 1.68 | 35 | 6 | $\mathbf{9 0 6}$ | 2.40 |  |
| 5832 | 648 | 104 | 5 | $\mathbf{7 5 8}$ | 1.17 | 136 | 6 | $\mathbf{1 0 0 1}$ | 1.54 |  |
| 10890 | 990 | 317 | 5 | $\mathbf{8 1 4}$ | 0.82 | 443 | 6 | $\mathbf{1 1 4 5}$ | 1.16 |  |
| 18252 | 1404 | 789 | 5 | $\mathbf{8 5 4}$ | 0.61 | 1122 | 6 | $\mathbf{1 2 3 2}$ | 0.88 |  |
| 28350 | 1890 | 1833 | 5 | $\mathbf{9 4 7}$ | 0.50 | 2222 | 6 | $\mathbf{1 1 6 9}$ | 0.62 |  |

Tab. 1. Contact problem with Coulomb friction.

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